## ILLC Project Course in Information Theory

Crash course
13 January - 17 January 2014
12:00 to 14:00
Student presentations
27 January - 31 January 2014
12:00 to 14:00

## Location

ILLC, room F1.15,
Science Park 107, Amsterdam

## Materials

informationtheory.weebly.com

## Contact

Mathias Winther Madsen
mathias.winther@gmail.com

Monday

Probability theory
Uncertainty and coding
Tuesday
The weak law of large numbers
The source coding theorem

## Wednesday

Random processes Arithmetic coding

## Thursday

Divergence
Kelly Gambling

## Friday

Kolmogorov Complexity
The limits of statistics

## Horse race

| Horse number | 1 | 2 | 3 | 4 |
| :--- | ---: | ---: | ---: | ---: |
| Win probability | .4 | .3 | .2 | .1 |
| Odds | 2 | 3 | 5 | 10 |


|  |  | Winner |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :--- |
|  |  |  | 1 | 2 | 3 |$⿻ 4$

## Horse race

| Horse number | 1 | 2 | 3 | 4 |
| :--- | ---: | ---: | ---: | ---: |
| Win probability | .4 | .3 | .2 | .1 |
| Odds | 2 | 3 | 5 | 10 |

Payoffs $R(b)=.4 \cdot 2 b+.3 \cdot 3 b+.2 \cdot 5 b+.1 \cdot 10 b$

$$
==p_{1} o_{1} b_{1}+p_{2} o_{2} b_{2}+\ldots+p_{n} o_{n} b_{n}
$$

## Degenerate Gambling



## Degenerate Gambling

Capital



## Degenerate Gambling

Rate of return:


## Degenerate Gambling

Rate of return:


## Optimal reinvestment



Daniel Bernoulli (1700-1782)


John Larry Kelly, Jr. (1923-1965)

## Optimal reinvestment

## The doubling rate:

$$
W==\log R==\log \frac{\text { Capital at time } i+1}{\text { Capital at time } i}
$$

$$
\text { So } 2^{W}=R \text {. }
$$

Long-run behaviour?

## Optimal reinvestment

## The doubling rate:

$$
W==\log R==\quad i p_{i} \log \left(b_{i} o_{i}\right)
$$

$$
\text { So } W=\quad \quad i p_{i} \log \left(o_{i}\right)-\quad i p_{i} \log \left(\frac{1}{b_{i}}\right)
$$

So that...?

## Optimal reinvestment

## Geometric expectation

$E[W]=\mathbf{p} \log$ bo
is maximized by propor-tional gambling ( $\mathbf{b}^{*}=\mathbf{p}$ ).

Arithmetic expectation

$$
E[R]=\mathbf{p b o}
$$

is maximized by degenerate gambling

## Optimal reinvestment



## Horse race

| Horse number | 1 | 2 | 3 | 4 |
| :--- | ---: | ---: | ---: | ---: |
| Win probability | .4 | .3 | .2 | .1 |
| Odds | 2 | 3 | 5 | 10 |

Payoffs

$$
\begin{aligned}
W(b)== & .4 \log (2 b)+.3 \log (3 b)+ \\
& .2 \log (5 b)+.1 \log (10 b)
\end{aligned}
$$

## Horse

 race| Horse number | 1 | 2 | 3 | 4 |
| :--- | ---: | ---: | ---: | ---: |
| Win probability | .4 | .3 | .2 | .1 |
| Odds | 2 | 3 | 5 | 10 |

Payoffs

$$
\begin{aligned}
W\left(b^{*}\right)== & .4 \log (2 \cdot .4)+.3 \log (3 \cdot .3)+ \\
& .2 \log (5 \cdot .2)+.1 \log (10 \cdot .1)
\end{aligned}
$$

## Arithmetic mean

$$
\frac{1}{n}\left(X_{1}+X_{2}+X_{3}+\ldots+X_{n}\right)
$$

Geometric mean

$$
\left(X_{1} \cdot X_{2} \cdot X_{3} \cdots X_{n}\right)^{1 / n}
$$

For example,

$$
\begin{gathered}
(1 / 2)(2+8)=5 \\
(2 \cdot 8)^{1 / 2}=4
\end{gathered}
$$

## The geometric mean

$$
\begin{gathered}
\left(X_{1} \cdot X_{2} \cdot X_{3} \cdot \cdots X_{n}\right)^{1 / n} \\
\text { is also equal to }
\end{gathered}
$$

$$
\left(2^{\log X_{1}+\log X_{2}+\log X_{3}+\ldots+\log X_{n}}\right)^{1 / n}
$$

which is equal to
$\left(2^{1 / n\left(\log X_{1}+\log X_{2}+\log X_{3}+\ldots+\log X_{n}\right)}\right)$
so for an ergodic process $X \ldots$

